

# A Refresher Exercise, Part 1

*Write predicate logic sentences for each of these facts:*

1. Joe, Sally, Bill, and Ellen are all members of the Elm St. Bridge Club.
2. Sally is married to Joe.
3. Bill is Ellen's brother.
4. The spouse of every married person in the club is also in the club.
5. The last meeting of the club was at Joe's house.

# One Possible Solution

## Objects

- Joe      Sally      Bill      Ellen      ElmStBridgeClub

## Predicates

- member( member, organization )
- married( spouse1, spouse2 )
- brother( person, theBrother )
- lastMeetingAt( organization, location )

## Functions

- houseOf( person )

## Sentences

1. member( Joe, ElmStBridgeClub )  
   member( Sally, ElmStBridgeClub )  
   member( Bill, ElmStBridgeClub )  
   member( Ellen, ElmStBridgeClub )
2. married( Sally, Joe )
3. brother( Ellen, Bill )
4. **forall** m, s  
   married( s, m ) **and**  
   memberOf( m, ElmStBridgeClub )  
   **implies** ( s, ElmStBridgeClub )
5. lastMeetingAt( ElmStBridgeClub, houseOf(Joe) )

## A Refresher Exercise, Part 2

*Given these facts:*

1. member( Joe, ElmStBridgeClub )  
member( Sally, ElmStBridgeClub )  
member( Bill, ElmStBridgeClub )  
member( Ellen, ElmStBridgeClub )
2. married( Sally, Joe )
3. brother( Ellen, Bill )
4. **forall** m, s  
married( s, m ) **and**  
memberOf( m, ElmStBridgeClub )  
**implies** ( s, ElmStBridgeClub )
5. lastMeetingAt( ElmStBridgeClub, houseOf(Joe) )

*Derive the following statements:*

1. The last meeting of the club was at Sally's house.
2. Ellen is not married.

*Feel free to define any commonsense axioms that you need in order to do your derivations.*

# One Possible Solution

## Commonsense Axioms

6. **forall** x, y  
    **married**( x, y ) **implies** **married**( y, x )
7. **forall** x, y  
    **married**( x, y ) **implies** **houseOf**( x ) = **houseOf**( y )
8. **forall** x, y  
    **married**( x, y ) **implies not** **brother**( x, y )  
    **forall** x, y  
    **married**( x, y ) **implies not** **brother**( y, x )
9. **forall** x, y, z  
    **married**( x, y ) **implies not** **married**( y, z )

## Derivations

# Where Are We?

We are exploring predicate logic with *modus ponens* as an inference rule.

- As a representation, logic gives a nice *formal language* in which to encode facts we believe to be true. Predicate logic allows us to represent a wide variety of facts about the world.
- As a reasoning mechanism, *modus ponens* lets us derive new facts in a straightforward, if cumbersome, way.

Using an inference rule to derive facts is a lot like using an operator to generate a successor to some state of the world.

In such a formulation, states are sets of sentences that describe a world, and an inference rule is an operator that allows you to generate successor states consisting of one more fact each.

## Where Are We?

So, we can evaluate an inference rule in much the same way as we would a search strategy.

*Modus ponens* is:

- completeness      Yes.
- optimality         It depends.
- time                 It depends, but ....
- space                It depends, but ....

Another interesting feature of logic is **soundness**. A sound inference rule derives only sentences that follow from the facts in the database. That is, if all sentences in a knowledge base are true, then it generates only true sentences.

*Modus ponens* is sound.

## Problems with this Approach

One difficulty with automating this sort of reasoning comes from the two kinds of sentences in our knowledge base:

- Hoosier( Eugene )
- Hoosier( x ) **implies** likes( x, Basketball )

Having two forms means that a program must be able to distinguish them and use them at the right times. This is an implementation problem.

Another difficulty with our current approach is that *modus ponens* only infers one fact from one fact and one implication. What if, for example, we wish to derive a new *implication*?

hoosier( x ) **implies** likes( x, Basketball )  
likes( x, Basketball ) **implies** likes( x, March )  
hoosier( x ) **implies** likes( x, March )

This is a problem with the logic itself.

We need a new sort of inference rule.

# Generalizing *Modus Ponens*

We can address both difficulties with one fix...

- **Represent facts as implications.**

Hoosier( x )

BECOMES

True **implies** Hoosier( x )

The semantics of **implies** make these sentences equivalent.

- **Modify *modus ponens* to work on implications.**

p  
p implies q  
q

BECOMES

True **implies** p<sub>i</sub>  
p<sub>1</sub> and ... p<sub>i</sub> and ... p<sub>n</sub> **implies** q  
p<sub>1</sub>... p<sub>i-1</sub> and p<sub>i+1</sub>... p<sub>n</sub> **implies** q



## Generalized *Modus Ponens*

This rule allows us to derive an implication...

True	<b>implies</b>	$p_i$
<u><math>p_1</math> and ... <math>p_i</math> and ... <math>p_n</math></u>	<b>implies</b>	$q$
$p_1 \dots p_{i-1}$ and $p_{i+1} \dots p_n$	<b>implies</b>	$q$

allows:

$a_1$ and ... $a_i$ and ... $a_n$	<b>implies</b>	$p_i$	
<u><math>p_1</math> and ... <math>p_i</math> and ... <math>p_n</math></u>	<b>implies</b>	$q$	$(p_i == a_i)$
$p_1 \dots p_{i-1}$ and $p_{i+1} \dots p_n$	and		
$a_1$ and ... $a_i$ and ... $a_n$	<b>implies</b>	$q$	

A benefit of this approach is that the reasoner now can have a single goal. To derive  $q$  from  $(p_1$  and ...  $p_i$  and ...  $p_n$  **implies**  $q$ ), use generalized modus ponens to

**drive**  $p_1$  and  $p_2$  and ...  $p_n$  **to** True

This is like recursion to a base case....